A step towards reconciling GALS industrial design with formal verification

Fatma Jebali

Join work with Frédéric Lang & Radu Mateescu

Inria Grenoble – France

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GALS: Globally Asynchronous, Locally Synchronous

Synchrony

- Several components governed by a single clock
- Instantaneous computations and communications
- Deterministic behaviour



```
1s, 2s, 3s, ...
```

Asynchrony

- Several subsystems executing at different speeds and interacting asynchronously
- Communications take arbitrary delays
- Nondeterministic behaviour





Em4: smart generation of PLCs (Programmable Logic Controllers)

- Software-based and multi-platform systems
- easy-to-use programming tool (block diagram model)
- application-oriented programming



Em4: smart generation of PLCs (Programmable Logic Controllers)





CADP <u>Construction and Analysis of Distributed Processes</u>

- Modular toolbox based on formal methods (nearly 50 tools up to now)
- Explicit-state verification
 - Model checking (µ-calculus, MCL)
 - Equivalence checking (bisimulations)
 - Visual checking
- Different techniques
 - Exhaustive
 - Partial
 - On the fly
 - Compositional
 - Distributed

http://cadp.inria.fr





The challenge

- Scalability of Model-checking (state space explosion)
- Expertise required to ensure correct and efficient specifications, write properties
- Automatic modeling and verification is recommended
 - Ability to handle industrial-size applications
 - Designer-friendly interfaces



Towards industrial-friendly formal verification



Towards industrial-friendly formal verification



Blocks: synchronous programs



Blocks as reactive systems

- Permanently interaction with the environment
- Constraints on one block
- Constraints on several blocks



Blocks as reactive systems



Blocks as asynchronously communicating systems

- Accurate design of
 - complex network topologies (bus, ring, star, etc.)
 - connection modes (point-to-point, multi-point, etc.)
 - communication protocol
- Existing languages
 - rigid topologies
 - point-to-point communications between separate synchronous systems



Blocks as asynchronously communicating systems

```
medium Med {receive Input : nat | send Output : nat} is
   perm Buffer : nat := 0
   select
     on Input -> if (Buffer == 0) then
                     Buffer := Input
                  else
                     null
                  end if
   []
      on ?Output -> if (Buffer == 0) then
                        Output := 0
                     else
                        Output := Buffer;
                        Buffer := 0
                     end if
   end select
end medium
```



Asynchronous composition and communication

- Define the system actors: blocks, environments, mediums
- Specify interactions between actors (message-passing synchronizations)
- Blocks execute arbitrarily and cyclically (active behaviour), triggering the activation of connected environments and mediums (passive behaviour)



Asynchronous composition and communication

```
system Main (in Switch : bool; in Sensor : nat;
             out Is On : bool; out Light : bool) is
  allocate Diagram1 as Controller1, Diagram2 as Controller2,
           Med as Medium1,
           Env Sensor as Envl
  temp Send Temperature : nat, Receive Temperature : nat
  network
     Controller1 (Switch; Sensor; ?Is On1) {?Send Temperature},
     Controller2 (?Light) {Receive Temperature}
  constrainedby
     Env1 (?Sensor)
  connectedby
     Medium1 {Send Temperature | ?Receive Temperature}
```

end system



This is just the tip of the iceberg...



First results

- F. Jebali, F. Lang, and R. Mateescu. GRL: A Specification Language for Globally Asynchronous Locally Synchronous Systems (Syntax and formal semantics). Research Report 8527, 82 pages, Inria, April 2014.
- F. Jebali, F. Lang, and R. Mateescu. GRL: A Specification Language for Globally Asynchronous Locally Synchronous Systems. In Proc. of ICFEM, October 2014.
- Automatic translator grl2int already designed and implemented, paper to be submitted (ETAPS 2015)

$$\begin{split} \rho_{const} &= \operatorname{init}(dl_{const}, \rho_{global})\\ \rho_{input} &= \bigoplus \operatorname{init}(dl_{in_i}, \rho_{global} \oplus \rho_{const}) \\ i \in 0.m \quad vars(al_{in_i}) \neq \epsilon \\ \{I_0\} \sigma.Ni, \rho_{var, \mu} \xrightarrow{vars(dl_{in_i})}_{i \quad \rho', \mu body} \\ \hline \{Ni \ (al_{in_0}| \dots | al_{in_m} | al_{out_0}| \dots | al_{out_n})\} \sigma, \rho, \mu \xrightarrow{vars(al_{in_i})}_{i \quad \rho', \mu body} \\ \hline \{Ni \ (al_{in_0}| \dots | al_{in_m} | al_{out_0}| \dots | al_{out_n})\} \sigma, \rho, \mu \xrightarrow{vars(al_{in_i})}_{i \quad \rho', \mu body} \\ \rho_{var} &= \rho_{global} \oplus \rho_{const} \oplus \rho_{input} \oplus \rho_{arg}) \\ \rho_{var} &= \rho_{global} \oplus \rho_{const} \oplus \rho_{input} \oplus \rho_{perm} \oplus \rho_{temp} \\ \rho' &= \rho \oplus \bigoplus_{i \in 0..n} \operatorname{update}(al_{out_i}, \operatorname{vars}(dl_{out_i}), \rho_{body}) \\ \oplus \bigoplus_{i \in 0..q} \operatorname{update}(al_{send_i}, \operatorname{vars}(dl_{send_i}), \rho_{body}) \\ \mu' &= \mu_{body} \oplus [\sigma.Bi \leftarrow \rho_{body} \cdot \operatorname{vars}(dl_{perm})] \end{split}$$



Ongoing work

- Automatic generation of GRL code from em4 software applications (large-scale industrial case studies)
- Industrial-friendly property language

Future work

- Automatic C code generation from GRL
- Connect GRL to other industrial frameworks to handle more instances of GALS systems
- Connect GRL to other verification tools



Thank you



