# **GRL: a Formal Language for the Specification of GALS Systems**

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### **ICFEM 2014**



• A set of synchronous systems composed asynchronously



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- Synchronous systems
  - Several components, one common clock
  - Instantaneous computations and communications
  - Deterministic behaviour





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# **Formal Verification of GALS Systems**

### Problem:

- Hard to design and debug
- Safety-critical applications

### Formal modeling and verification:

- Powerful automatic tools
- Correctness of the design process

### However:

- Expertise in formal methods required
- Scalability to industrial-size applications

### Solution:

### **GRL (GALS Representation Language)**



# Rationale for GRL (GALS Representation Language)

### • User convenience

- Unified language (synchronous and asynchronous)
- Modular modeling
- Abstraction
- Easy-to-use
- Efficient formal verification
  - Formal semantics
  - Pivot language (industrial tools, CADP [1] toolbox)

[1] Construction and Analysis of Distributed Processes http://cadp.inria.fr/



### **GRL in a nutshell**

- Synchronous systems
  - Blocks: synchronous behaviour
  - Based on the dataflow model
- Asynchronous composition
  - Mediums: communication between blocks
  - Environments: external constraints
  - Inspired by process algebraic languages
- Imperative flavour





### **Running Example**





# **Blocks**

- Cyclic behaviour (active):
  - Discrete deterministic steps
    - 1. Consume inputs
    - 2. Compute a reaction
    - 3. Produce outputs
  - Memory maintained: permanent variables
  - Atomic
- Composition of subblocks







# **Blocks**

- Cyclic behaviour (active):
  - Discrete deterministic steps
    - 1. Consume inputs
    - 2. Compute a reaction
    - 3. Produce outputs
  - Memory maintained: permanent variables
  - Atomic
- Composition of subblocks
- Receive, Send: asynchronous communication





block Heater (in Switch : bool; in Sensor : nat; out Is\_On : bool) is

allocate Comparator [Strictly\_Inferior] as B02, NUM [3] as B03, AND as B04,

temp c1 : bool, c2 : nat

B03 (?c2); B02 (\_; Sensor, c2; ?c1); B04 (Switch; c1; ?ls\_On)

end block





Physical interactions

**block** Heater (in Switch : **bool**; in Sensor : **nat**; **out** Is\_On : **bool**) is

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allocate Comparator [Strictly\_Inferior] as B02, NUM [3] as B03, AND as B04, Creation of instances Separate memories

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end block



Temporary

variables



**block** Heater (in Switch : **bool**; in Sensor : **nat**; **out** Is\_On : **bool**) is

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block Heater (in Switch : bool; in Sensor : nat; out Is\_On : bool) is

























- Modeling of asynchronous communication
- Activated on demand (passive)
  - Several connected blocks, different instants
  - Nondeterminism
- Signal statements to control activation





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medium Coord {receive apo : nat | send lock; up; down : bool |

receive lp, up, dp : bool | send app : nat |

receive ls, us, ds : bool | send aps : nat} is

Buffers for transited data

perm lock\_bu : bool := true, up\_bu, down\_bu : bool := false, apo\_bu : nat := 0

#### select

```
on lp, up, dp -> lock_bu := lp; up_bu := up; down_bu := dp
```

- [] **on** ls, us, ds **->** lock\_bu **:=** ls; up\_bu **:=** us; down\_bu **:=** ds
- [] on apo -> apo\_bu := apo
- [] **on** ?app -> app := apo\_bu
- [] on ?aps -> aps := apo\_bu

```
[] on ?lock,?up, ?down -> lock := lock_bu ; up := up_bu ; down := down_bu
```

end select

#### end medium



medium Coord {receive apo : nat | send lock; up; down : bool |
 receive lp, up, dp : bool | send app : nat |
 receive ls, us, ds : bool | send aps : nat} is



#### end medium

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- [] on apo -> apo\_bu := apo
- [] **on** ?app -> app := apo\_bu
- [] on ?aps -> aps := apo\_bu

[] on ?lock,?up, ?down -> lock := lock\_bu ; up := up\_bu ; down := down\_bu
end select

#### end medium





#### end medium



### **Environments**

### Modeling of constraints

- Logical constraints between blocks
- Physical constraints
- Activated on demand (passive)
- Signal statements to control activation



### **Environments: FCS Example**

environment Conc (out p\_tok:bool | out

Safety state of Prim and Sec

bool) <mark>is</mark>

perm p\_alive, s\_alive:bool := true

if p\_alive then

select

on ?p\_tok -> p\_tok := true -- primary responds

[] p\_alive := false -- primary fails

end select

elsif s\_alive then

select

on ?s\_tok -> s\_tok := true -- secondary responds

[] s\_alive := false -- secondary fails

end select

else

on ?alarm -> alarm := true

end if

end environment







### **Systems**

- Composition of blocks, mediums, and environments
- No direct connection between blocks
- Communication between blocks and mediums (resp., environments) by message-passing rendezvous





system FlightControlSystem (p\_ord, s\_ord : nat, alarm : bool) is

allocate FBWComp as Prim, FBWCom as Sec, Ail as Ail, Alarmer as Alarmer,

Conc as Conc, Ctrl [10] as Ctrl, Coord as Coord

**temp** p\_tok : **bool**, p\_pos: **nat**, p\_lck, p\_up, p\_dwn : **bool**,

s\_tok : **bool**, **s**\_pos: **nat**, s\_lck, s\_up, s\_dwn : **bool**,

c\_pos, pos : nat, lck, up, dwn : bool, safe, ok: bool

#### network

Prim (p\_tok; p\_ord) {p\_pos; ?p\_lck, ?p\_up, ?p\_dwn},

```
Sec (s_tok; s_ord) {s_pos; ?s_lck, ?s_up, ?s_dwn},
```

```
Ail (ok; ?c_pos) {lck, up, dwn; ?pos},
```

Alarmer (safe; ?alarm)

#### constrainedby

Conc (?p\_tok | ?s\_tok | ?safe),

Ctrl (c\_pos | ?ok)

#### connectedby

Coord {pos | ?lck, ?up, ?dwn | p\_lck, p\_up, p\_dwn | ?p\_pos | s\_lck, s\_up, s\_dwn | ?s\_pos}

#### end system

Creation of instances

**system** FlightControlSystem (p\_ord, s\_ord : **nat**, alarm : **bool**) is allocate FBWComp as Prim, FBWCom as Sec, Ail as Ail, Alarmer as Alarmer, Conc as Conc, Ctrl [10] as Ctrl, Coord as Coord **temp** p\_tok : **bool**, p\_pos: **nat**, p\_lck, p\_up, p\_dwn : **bool**, s\_tok : **bool**, **s**\_pos: **nat**, s\_lck, s\_up, s\_dwn : **bool**, c\_pos, pos : nat, lck, up, dwn : bool, safe, ok: bool network Prim (p\_tok; p\_ord) {p\_pos; ?p\_lck, ?p\_up, ?p\_dwn}, Block Sec (s\_tok; s\_ord) {s\_pos; ?s\_lck, ?s\_up, ?s\_dwn}, invocations Ail (ok; ?c\_pos) {lck, up, dwn; ?pos}, Alarmer (safe; ?alarm)

#### constrainedby

Conc (?p\_tok | ?s\_tok | ?safe),

Ctrl (c\_pos | ?ok)

#### connectedby

Coord {pos | ?lck, ?up, ?dwn | p\_lck, p\_up, p\_dwn | ?p\_pos | s\_lck, s\_up, s\_dwn | ?s\_pos}

#### end system



Coord {pos | ?lck, ?up, ?dwn | p\_lck, p\_up, p\_dwn | ?p\_pos | s\_lck, s\_up, s\_dwn | ?s\_pos}

#### end system



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#### network

Prim (p\_tok; p\_ord) {p\_pos; ?p\_lck, ?p\_up, ?p\_dwn}, Sec (s\_tok; s\_ord) {s\_pos; ?s\_lck, ?s\_up, ?s\_dwn}, Ail (ok; ?c\_pos) {lck, up, dwn; ?pos}, Alarmer (safe; ?alarm)

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#### end system

Visible from the



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## **Formal Semantics of GRL**

- Labelled transition systems
  - States: union of the memories of blocks, mediums, and environments
  - Initial state: initial values of memories
  - Labels: execution of blocks
    - + visible inputs/outputs + visible receives/sends
  - Transition function: atomic execution of blocks with connected mediums and environments
- 145 rules of static semantics [2]
- 24 rules of structural operational semantics [2]

[2] available in a technical report of 130 pages

### **Tools for GRL**

- GRL2LNT(20,000 lines):
  - Parser (2,000 lines): lexical and syntactic analysis
  - Automated translator to LNT, input language of CADP
  - Accurate and concise LNT
  - Improve scalability of model checking
- Enabled access to CADP
  - More than 40 tools
  - Explicit state exploration
  - Model-checking, equivalence checking, visual checking



## **Results for the FCS Example**

- State space generation
  - 2,653 states
  - 7,406 transitions
- Reduction with branching bisimulation
  - 5 states
  - 1,287 transitions
- Formal verification with CADP enabled



# **Conclusion: GRL**

- Versatile and modular description of
  - Synchronous systems
  - Asynchronous communication
  - Environment constraints
- Expressive and general-purpose
- Close to graphical data flow used in industry
- Easier to learn than full-fledged process algebra
- Efficient verification with CADP



# **Conclusion: ongoing work**

- GRL and GRL2LNT applied on an industrial project
  - Crouzet Automation (Schneider Electric)
  - Networks of Programmable Logic Controllers

### ➔ Positive feedback

- Development of off-the-shelf blocks, mediums, and environments
- Automated GRL generation from industrial tools
  - ➔ Automated verification chain
- Connection to synchronous verification tools (future work)

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### **Thank You**

